Imperial College London



Teechain: A Secure Payment Network with Asynchronous Blockchain Access

Joshua Lind

Florian Kelbert

Oded Naor

Emin Gün Sirer

Ittay Eyal

Peter Pietzuch



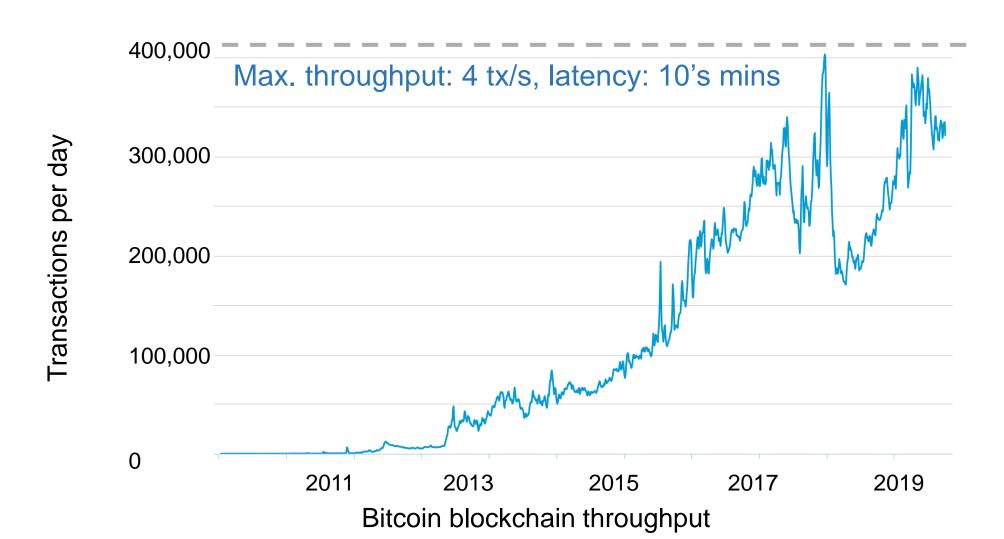
joshua.lind11@imperial.ac.uk Imperial College London





Blockchains aren't scaling!

Consensus is slow: all nodes must agree on all transactions!



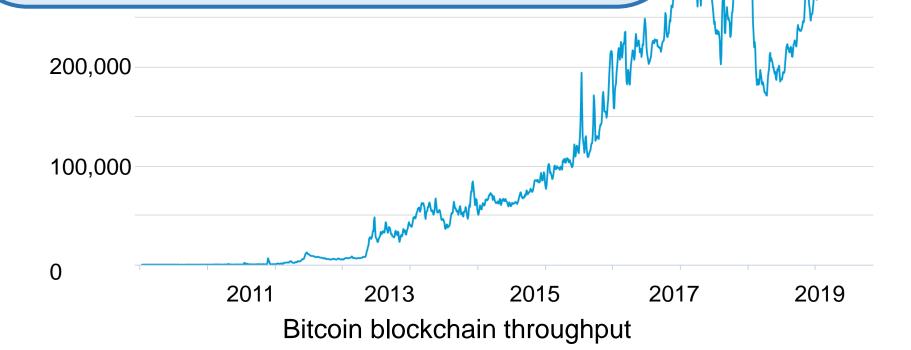


Transactions per

Consensus isn't scaling!

How can we reach the throughput of Visa, Mastercard or Paypal?

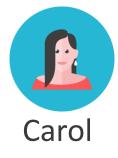
100,000 tx/s





Execute payments off-chain!

- Parties pay each other directly









Blockchain



Execute payments off-chain!

Parties pay each other directly

Payment channels:

- Point to point payments

Carol

Dave

Blockchain

Bob



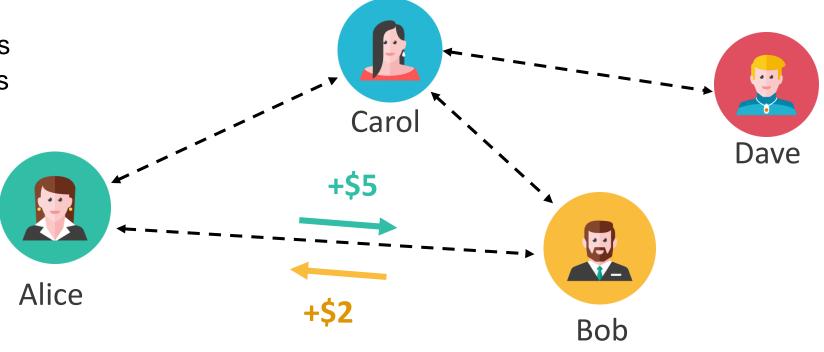
Execute payments off-chain!

- Parties pay each other directly

Payment channels:

Point to point payments

Bi-directional payments



Blockchain

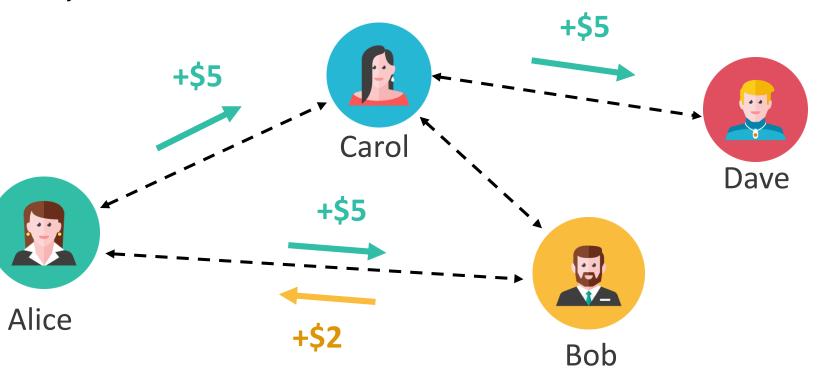


Execute payments off-chain!

Parties pay each other directly

Payment channels:

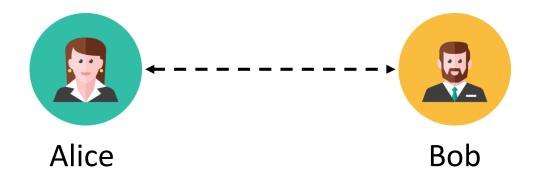
- Point to point payments
- Bi-directional payments
- Multi-hop payments



Blockchain



How do payment channels work?

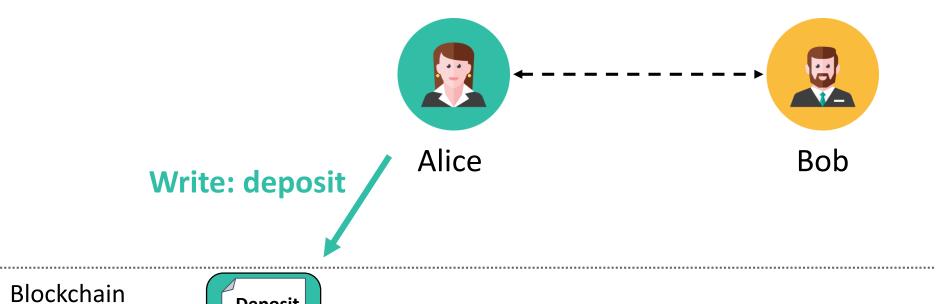




How do payment channels work?

- 3 phases: **setup**, payments, settlement

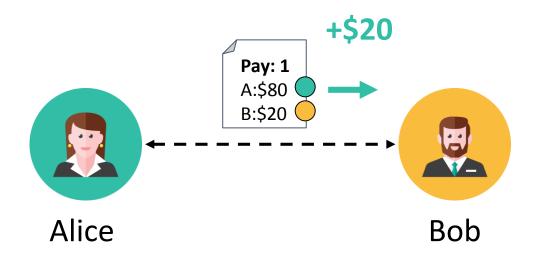
Deposit

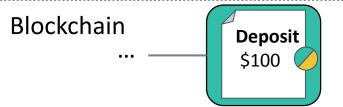


9



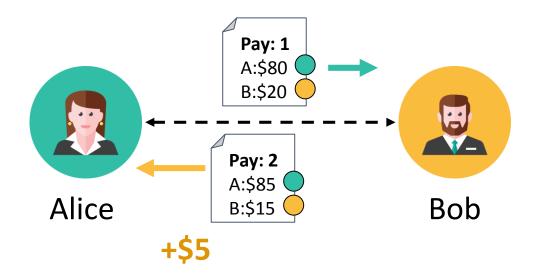
How do payment channels work?

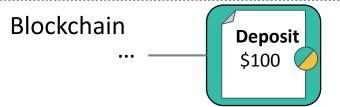






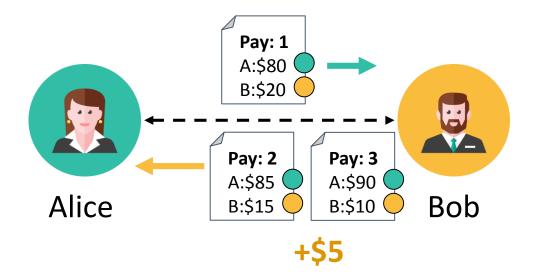
How do payment channels work?

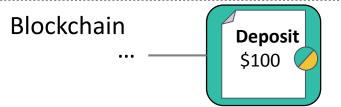






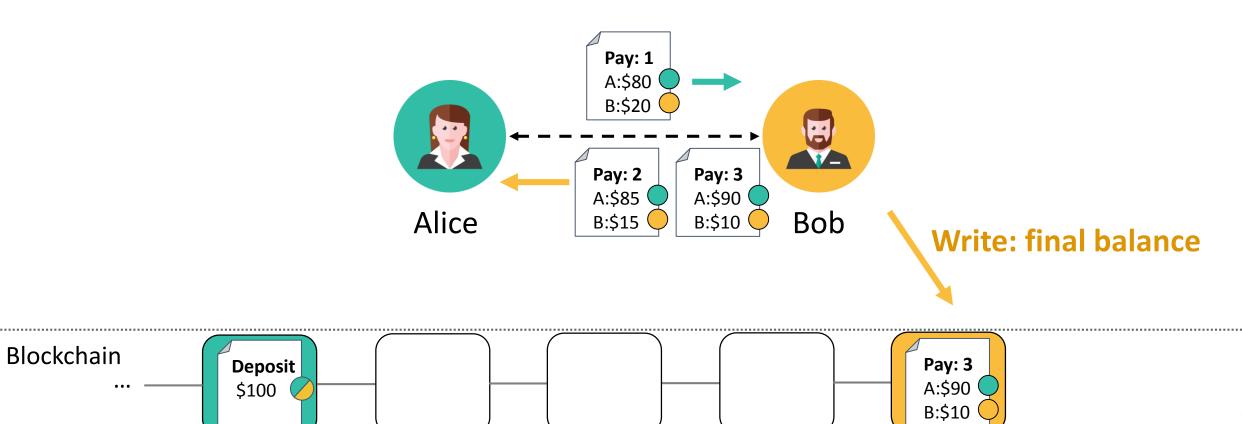
How do payment channels work?







How do payment channels work?

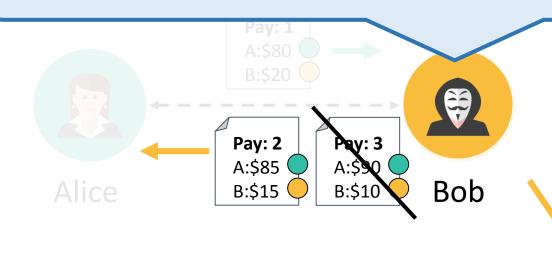




Roll-back attacks!

How do payr
- 3 phases: se

What if Bob misbehaves and writes an old balance to the blockchain?



Write: old balance

Blockchain

Deposit
\$100

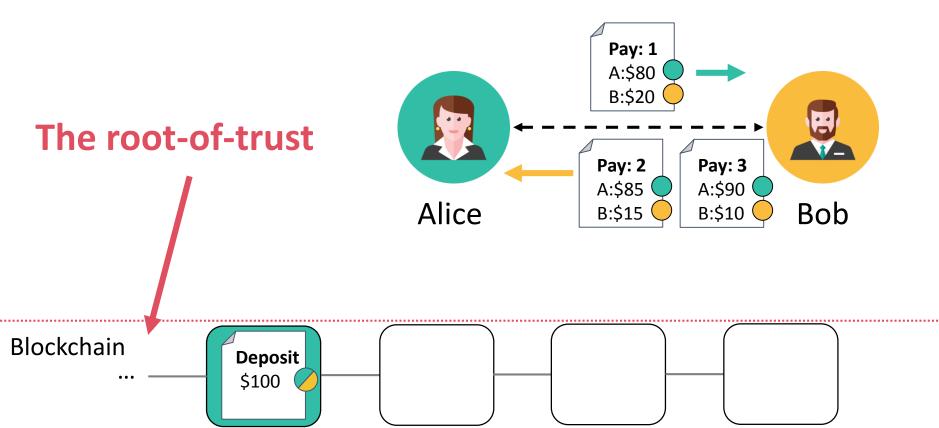
Pay: 2

A:\$85

B:\$15

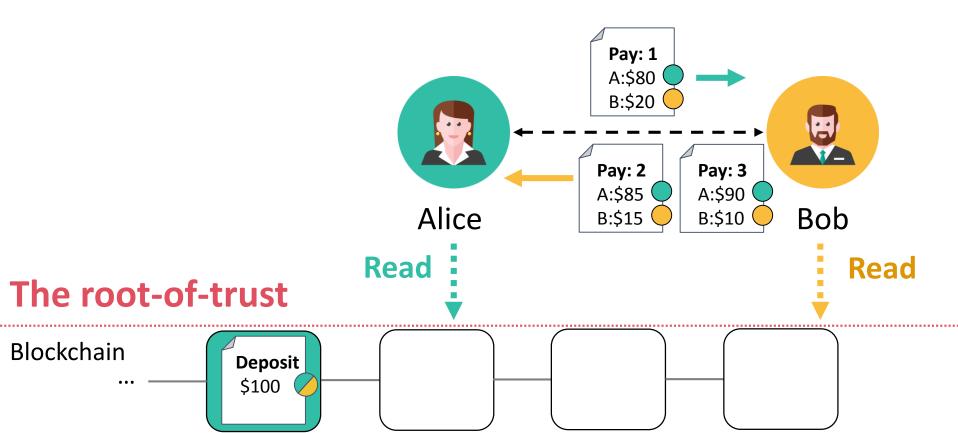


- Monitor the blockchain (root-of-trust)
- React within reaction time (Δ)
- Final balance on the blockchain



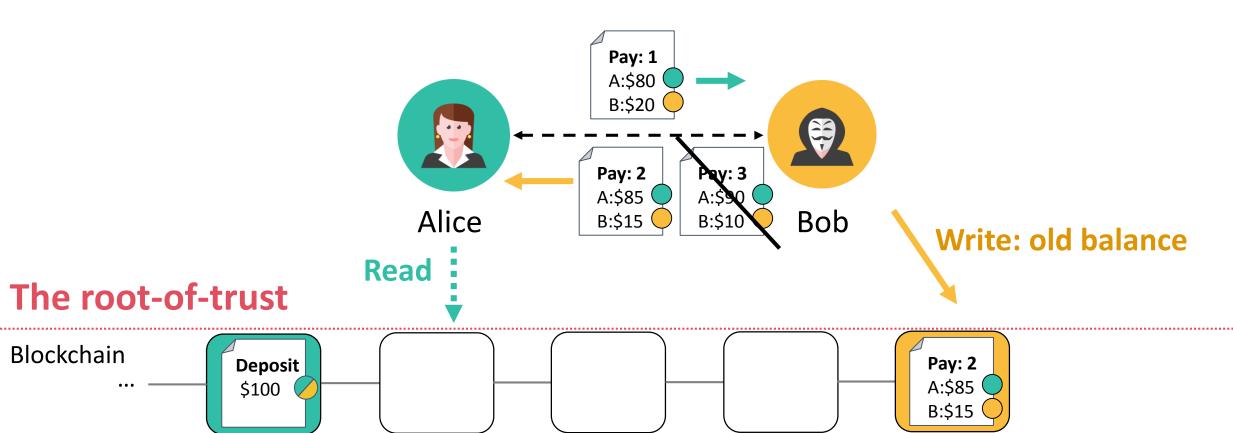


- Monitor the blockchain (root-of-trust)
- React within reaction time (Δ)
- Final balance on the blockchain





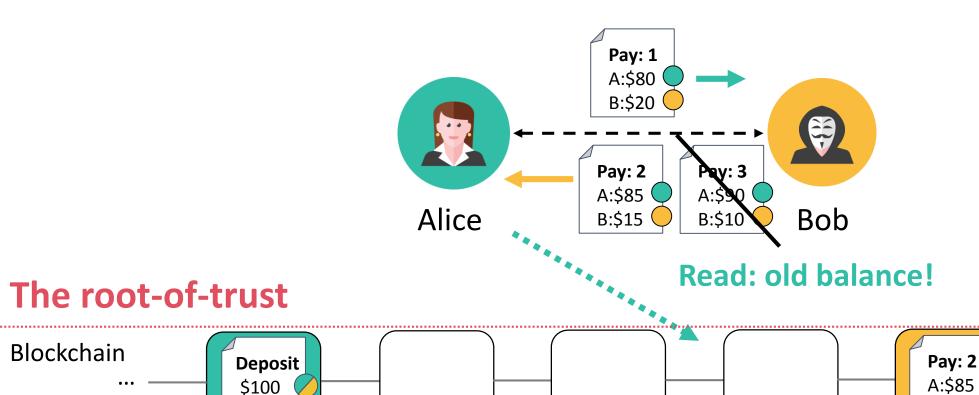
- Monitor the blockchain (root-of-trust)
- React within reaction time (Δ)
- Final balance on the blockchain





Existing solutions to roll-back attacks:

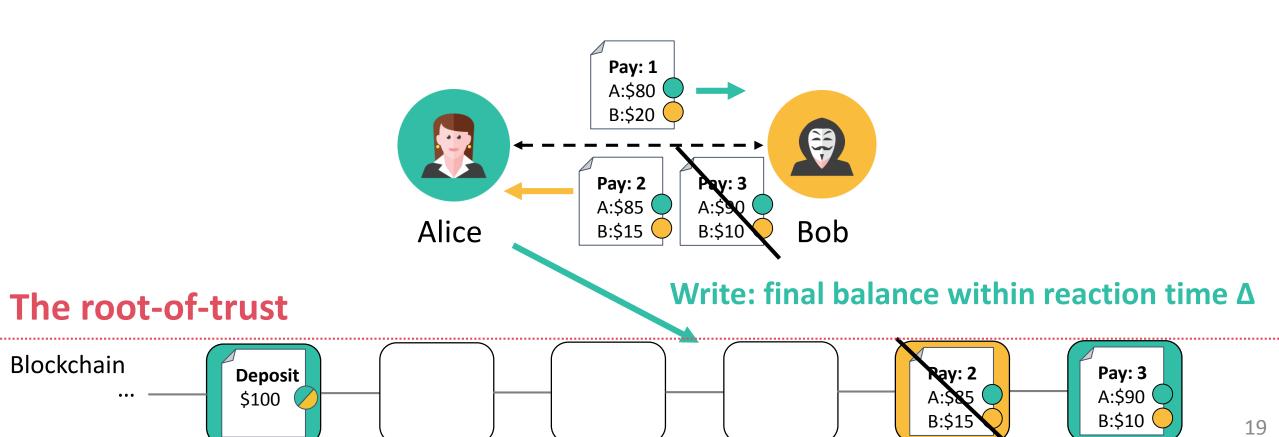
- Monitor the blockchain (root-of-trust)
- React within reaction time (Δ)
- Final balance on the blockchain



B:\$15



- Monitor the blockchain (root-of-trust)
- React within reaction time (Δ)
- Final balance on the blockchain





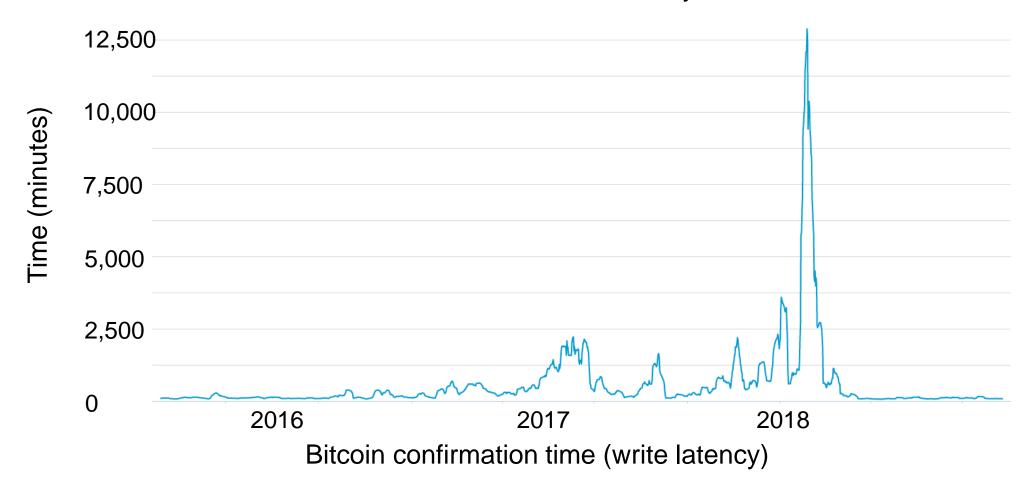
Reaction times (Δ) require synchronous blockchain access:

- Assume: parties can read/write within Δ
- But: blockchains are best-effort. No read/write latency bounds!



Reaction times (Δ) require synchronous blockchain access:

- Assume: parties can read/write within Δ
- But: blockchains are best-effort. No read/write latency bounds!

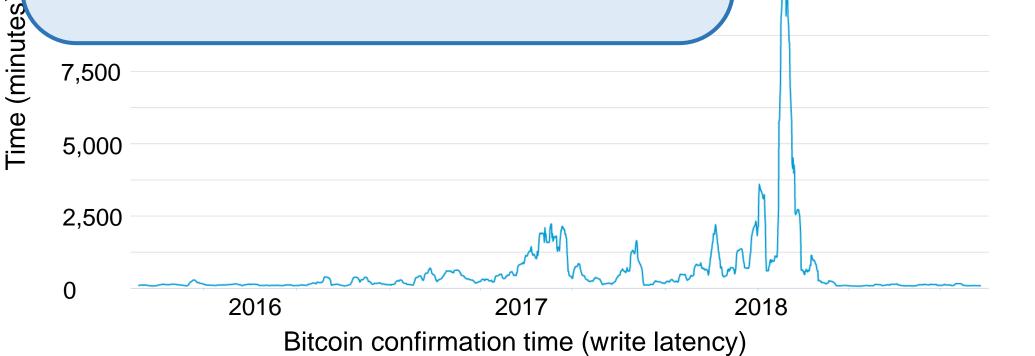




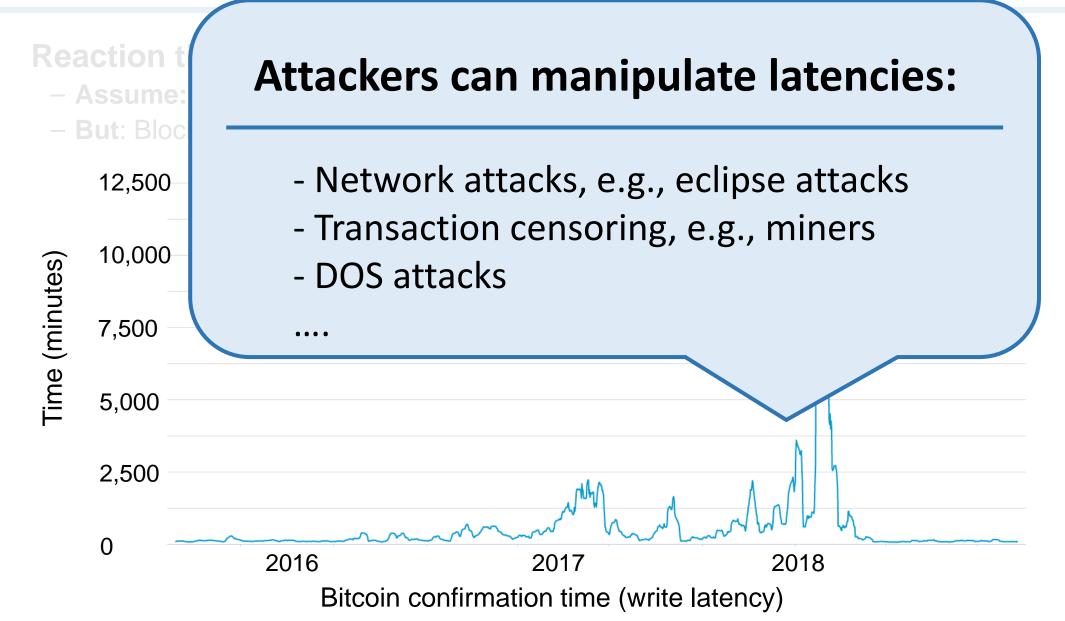
Access:

Spam/Congestion attack!

Transactions took > **7 days** to be written to the blockchain!







I-back attacks What value for reaction time (Δ) ? **Trade-off: Large Δ (weeks):** hard to attack, slow fund access Small Δ (mins): easy to attack, quick fund access Write: final balance within **\Delta** Pay: 3 A:\$90 B:\$10 2016 2017 2018 Bitcoin confirmation time (write latency)



Teechain: Challenges and roadmap

Asynchronous blockchain access (no read/write latency bounds):

Challenge 1: removing the blockchain as root-of-trust (RoT)

Idea: treasury as new RoT for payments



Challenge 2: realizing treasuries for blockchains

Idea: decentralized *treasury committees*

Idea: *trusted execution* to secure committees

Challenge 3: consensus in treasury committees

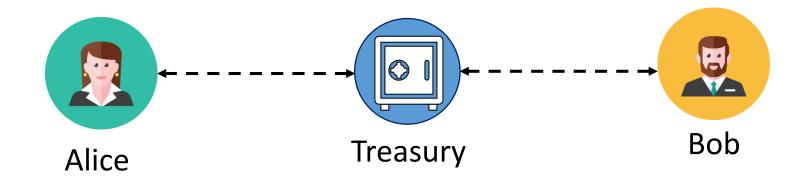
Idea: force-freeze chain replication





Introduce another root-of-trust (RoT): treasury

- Controls funds, balances and payments
- Prevents misbehaviour
- Only settle channels once → prevents roll-backs!



Blockchain

••

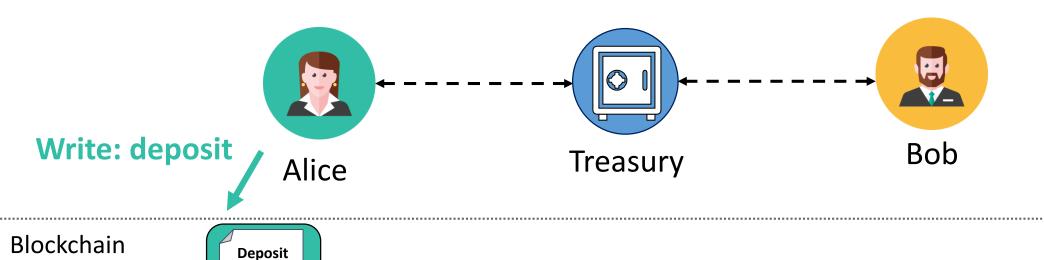


Introduce another root-of-trust (RoT): treasury

- Controls funds, balances and payments
- Prevents misbehaviour

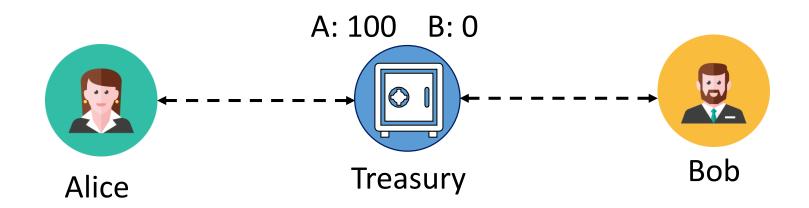
\$100

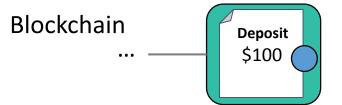
— Only settle channels once → prevents roll-backs!





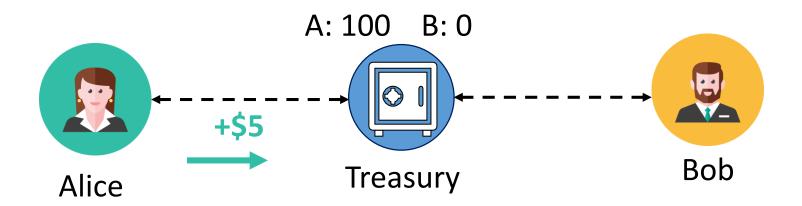
- Controls funds, balances and payments
- Prevents misbehaviour
- Only settle channels once → prevents roll-backs!

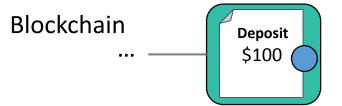






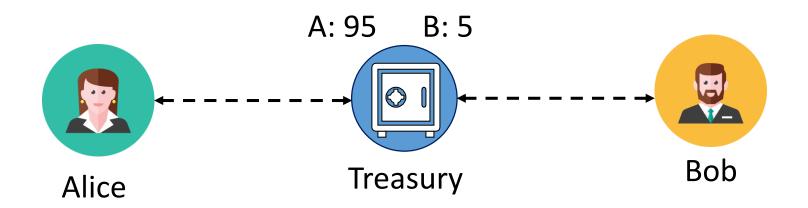
- Controls funds, balances and payments
- Prevents misbehaviour
- Only settle channels once → prevents roll-backs!

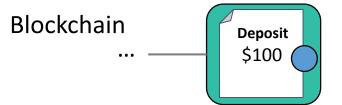






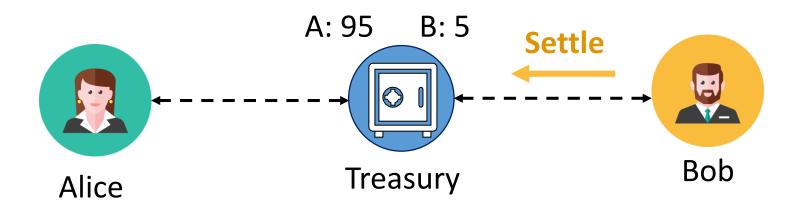
- Controls funds, balances and payments
- Prevents misbehaviour
- Only settle channels once → prevents roll-backs!

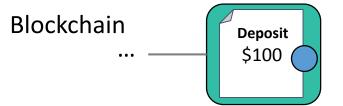






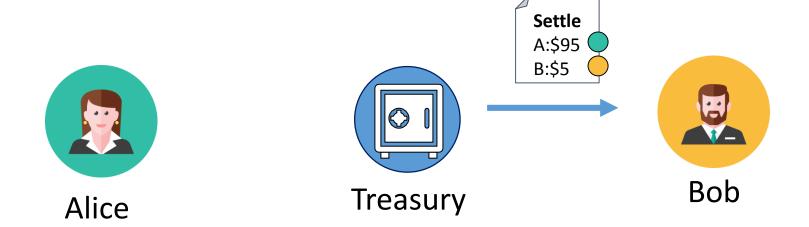
- Controls funds, balances and payments
- Prevents misbehaviour
- Only settle channels once → prevents roll-backs!





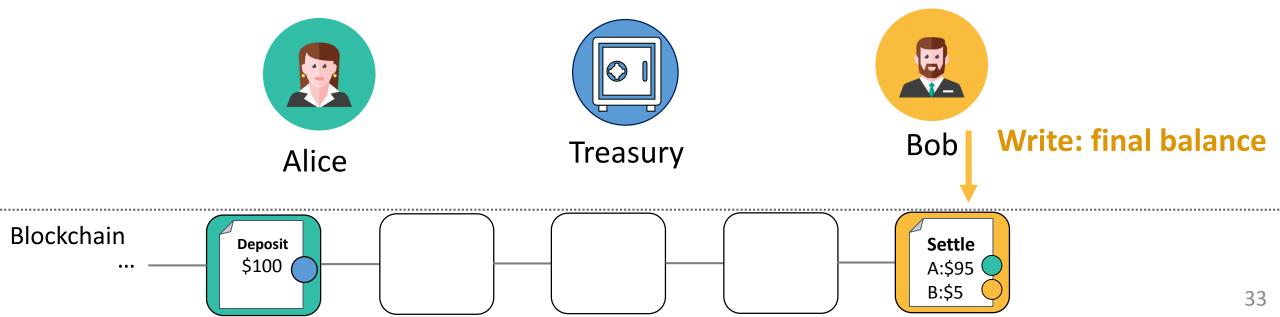


- Controls funds, balances and payments
- Prevents misbehaviour
- Only settle channels once → prevents roll-backs!





- Controls funds, balances and payments
- Prevents misbehaviour
- Only settle channels once → prevents roll-backs!





Treasury prevents roll-backs!

No roll-backs → no reaction times:

<u>Asynchronous blockchain access</u>







Treasury



Bob

Write: final balance

- Unbounded latency!

Blockchain

Deposit \$100







Challenge 1: Pomoving the blockshoin of RoT

Introduce another roo

- Controls funds, balan
- Prevents misbehavior
- Only settle channels

Can we realize treasuries?

How do we realize treasuries for blockchains?



Alice



Treasury

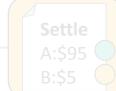


Bob

Write: Final Balance

Blockchain

Deposit \$100





Challenge 2: Realizing treasuries for blockchains

Design treasury to:

- Avoid absolute trust (parties are selfish!)
- Avoid centralization
- Integrate with most blockchains (e.g. no smart contracts!)



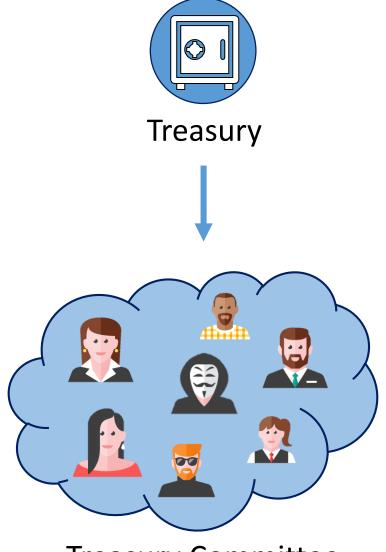


Design treasury to:

- Avoid absolute trust (parties are selfish!)
- Avoid centralization
- Integrate with most blockchains (e.g. no smart contracts!)

Use a committee!

- General solution: well studied for blockchains
- Decentralized: distribute trust
- Fault tolerant: crash and Byzantine failures





Treasury committee:

- Choose **n** parties in the network
- Require m parties to agree before accessing funds
- Use m-out-of-n transactions



Blockchain

••

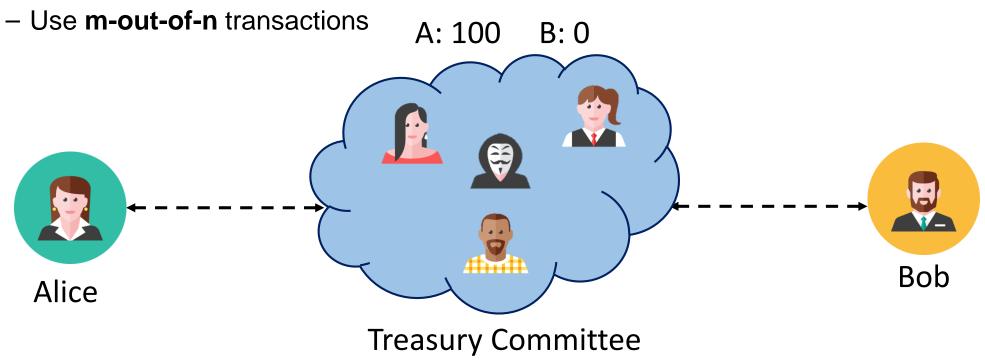


- Choose **n** parties in the network
- Require m parties to agree before accessing funds
- Use m-out-of-n transactions



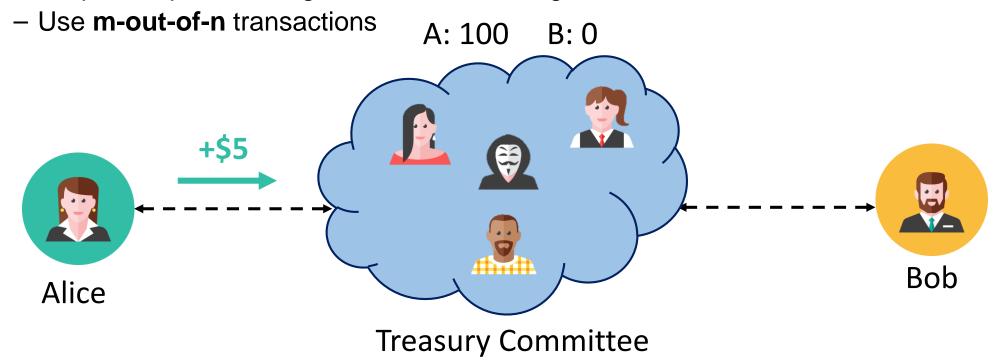


- Choose **n** parties in the network
- Require m parties to agree before accessing funds



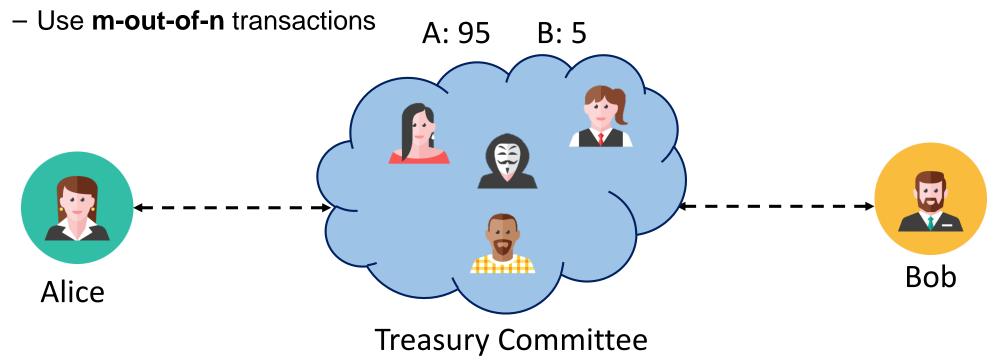


- Choose **n** parties in the network
- Require m parties to agree before accessing funds





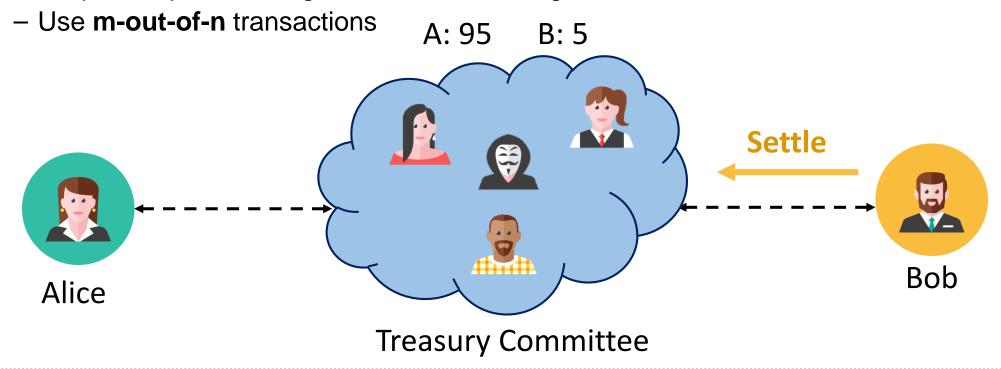
- Choose **n** parties in the network
- Require m parties to agree before accessing funds







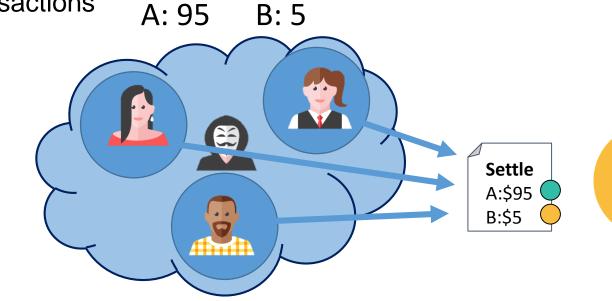
- Choose **n** parties in the network
- Require **m** parties to agree before accessing funds





Treasury committee:

- Choose **n** parties in the network
- Require m parties to agree before accessing funds
- Use m-out-of-n transactions



Treasury Committee

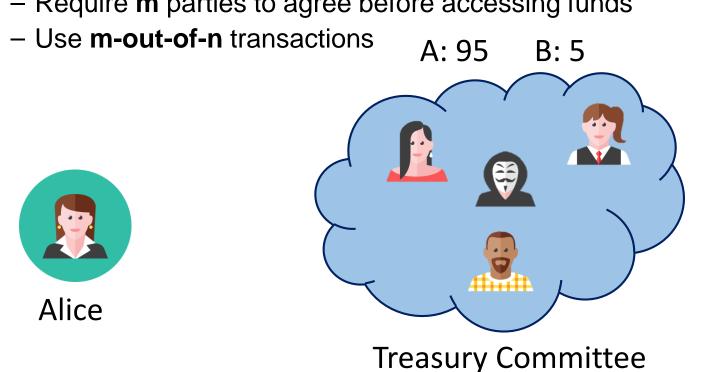


Blockchain 3-out-of-4 Deposit \$100

Bob



- Choose **n** parties in the network
- Require m parties to agree before accessing funds







Challenge 2

chains

Trust is distributed!

Treasury commit

- Choose n parties
- Require m parties
- Use m-out-of-n t

m treasuries must collude together to steal the deposit!



Treasury Committee

Treasury committee size?

How large should **m** and **n** be?



Blockchain

3-out-of-4 Deposit \$100 **Settle**A:\$95
B:\$5

ries for blockchains



Existing solutions:

- Large committees for security: e.g. Elastico, Algorand...
- But this is difficult at scale! (consensus..)





Existing solutions:

- Large committees for security: e.g. Elastico, Algorand...
- But this is difficult at scale! (consensus..)



Smaller committees? Use trusted execution!

- Confidentiality + integrity guarantees
- Only trust hardware and manufacturer (don't trust people!)

Many trusted execution environments (TEEs):

- Commodity: Intel SGX, ARM TrustZone, AMD SEV...
- Up-and-coming: KeyStone Enclave, Multizone, OP-TEE, Sanctum...













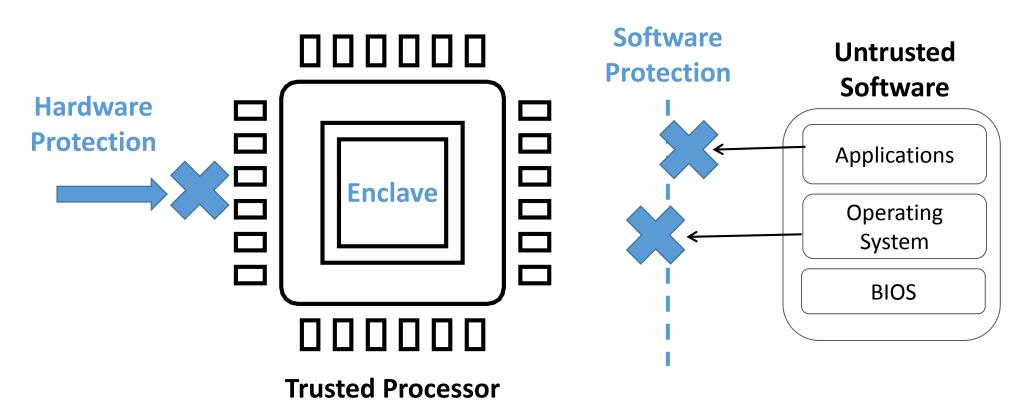


Background: Intel software guard extensions (SGX)

Intel SGX provides confidentiality and integrity for enclaves:

- Software protection: OS, BIOS, other applications
- Physical attacks: DRAM, Disk, System Bus







2-out-of-3

Deposit

\$10

Use TEEs to secure committee members

Increase attack costs: reduce committee size

TEEs are not silver bullets:

- Existing attacks: e.g. Foreshadow [USENIX SEC'18]
- Combine TEEs + committees: defence-in-depth

"Configurable security" per deposit:

- Parties decide m-out-of-n: no "one size fits all"
- TEE heterogeneity: avoid centralization/attacks
- Weigh-up deposit risk: e.g.,
 - **\$10:** 2-out-of-3 committee
 - **\$100:** 3-out-of-4 committee



Treasury Committee





How do we maintain treasury agreement?

Peer-to-peer network → not fully connected (e.g. NATs and firewalls)



How do we maintain treasury agreement?

Peer-to-peer network → not fully connected (e.g. NATs and firewalls)

Use chain replication:

- Strong consistency: using a chain topology
- Efficient: update in O(n) messages
- Easy to reason about: avoid bugs!





How do we maintain treasury agreement?

Peer-to-peer network → not fully connected (e.g. NATs and firewalls)

Use chain replication:

- Strong consistency: using a chain topology
- Efficient: update in O(n) messages
- Easy to reason about: avoid bugs!





How do we maintain treasury agreement?

Peer-to-peer network → not fully connected (e.g. NATs and firewalls)

Use chain replication:

- Strong consistency: using a chain topology
- Efficient: update in O(n) messages
- Easy to reason about: avoid bugs!





What about failures?

Failures allow roll-back/replay attacks: Introduce force-freeze chain replication (see the paper!)

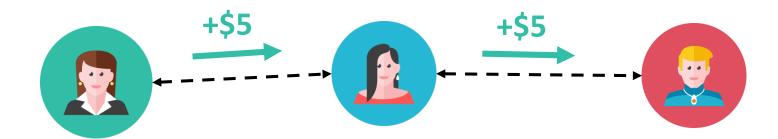




See the paper!

Multi-hop protocol

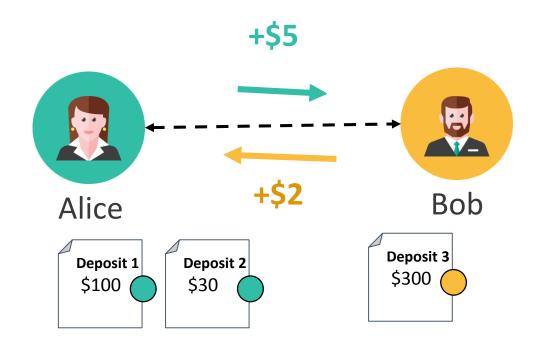
- Multi-phase commit



Dynamic fund deposits

- Add/remove funds dynamically!

More features/optimizations!





Teechain: Implementation

Teechain Network:

- Bitcoin BTC blockchain (ported Bitcoin core)
- Intel SGX (20k C++ LoC inside TEE)
- 65k untrusted C++ LoC

Open-source (available and functional badges)

- Github: https://github.com/lsds/Teechain
- **Visit us**: teechain.network







Teechain: Evaluation

Evaluation questions:

- 1. How well do **payment channels** perform?
- 2. How well do **multi-hop payments** perform?
- 3. Does Teechain scale out?

Baseline comparison:

- State of the art Lightning Network for Bitcoin
- Requires synchronous blockchain access

Experimental setup:

- 35 SGX machines across London, New York and Haifa
- Intel Xeon E3-1280 v5 32GB RAM





Payment channel: London -- New York

- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)



Payment channel: London -- New York

- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)

Maximum throughput

Throughput (tx/s)

20000

20000

150000

0



Payment channel: London -- New York

- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)

Maximum throughput



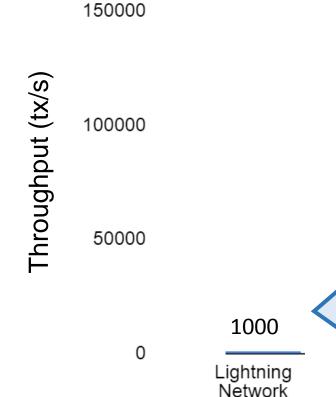
0 Lightning Network



Payment channel: London -- New York

- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)

Maximum throughput



Limited throughput

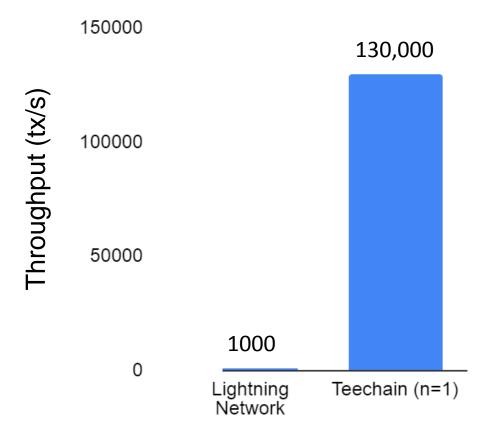
Each payment requires multiple message exchanges



Payment channel: London -- New York

- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)

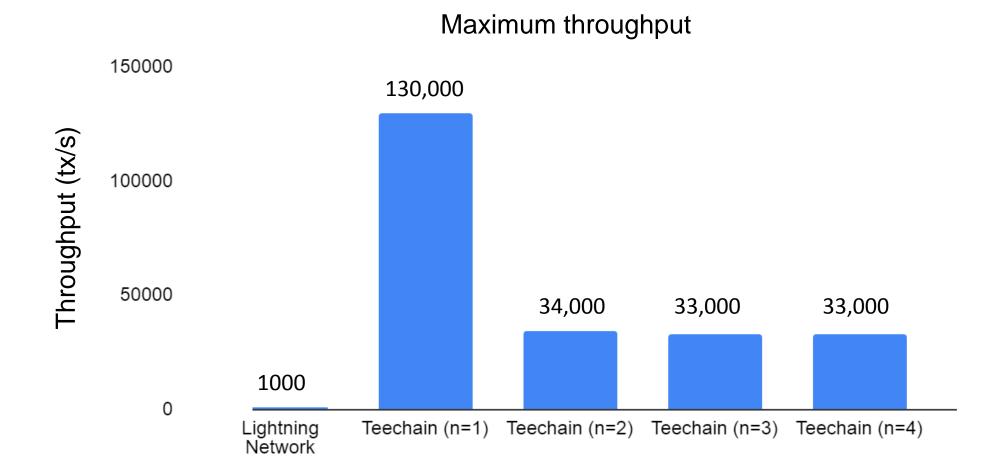




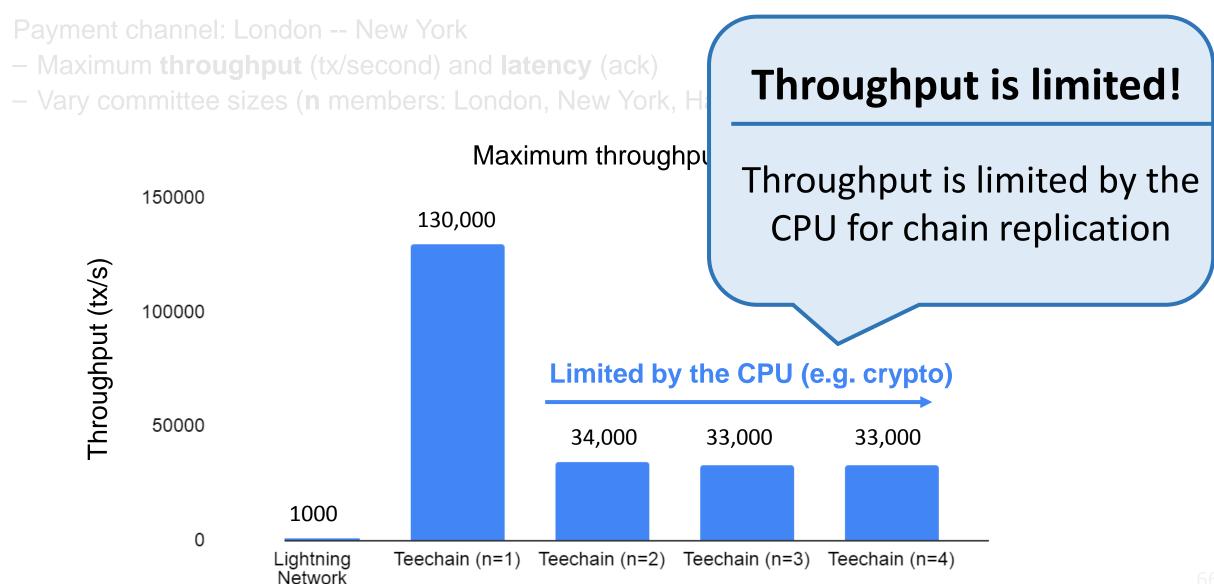


Payment channel: London -- New York

- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)









Payment channel: London -- New York

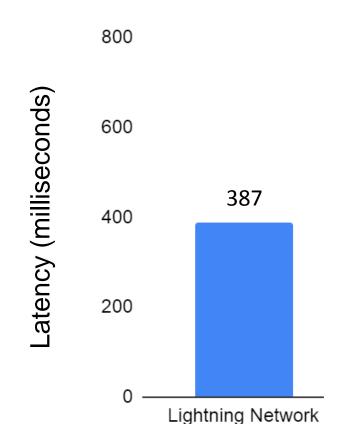
- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)





Payment channel: London -- New York

- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)



Maximum latency

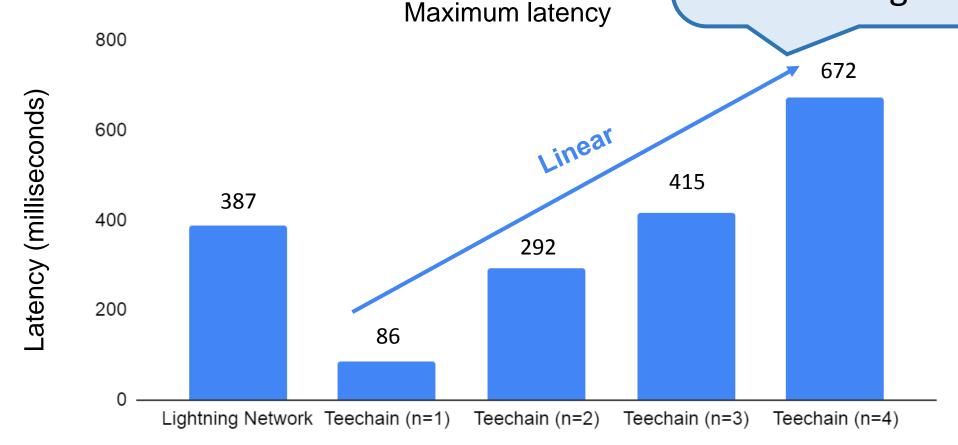


Payment channel: London -- New York

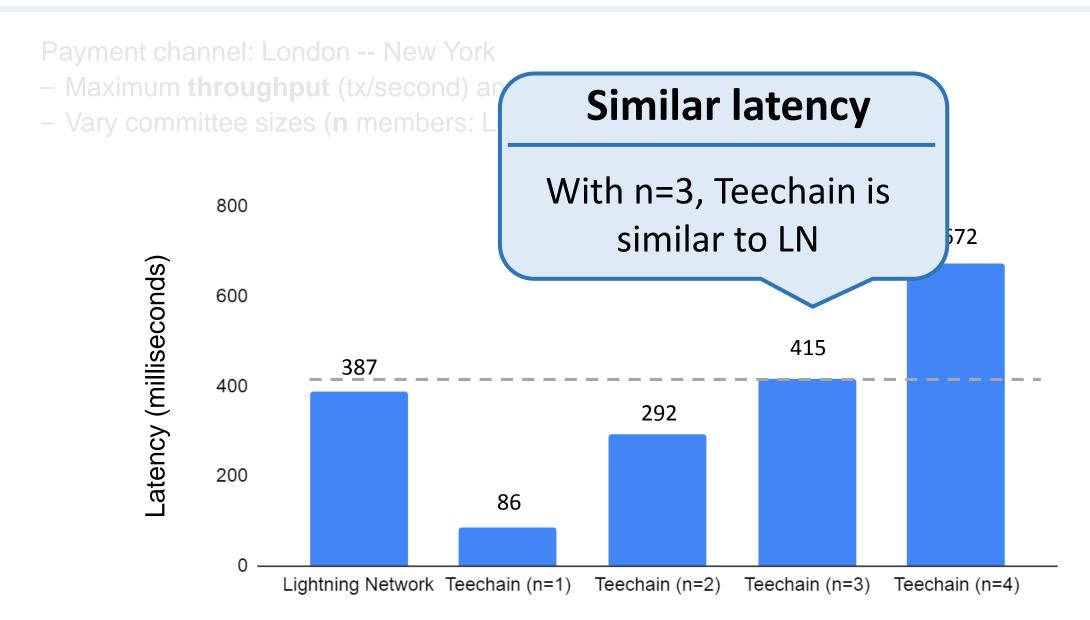
- Maximum throughput (tx/second) and latency (ack)
- Vary committee sizes (n members: London, New York, Haifa)

Latency grows

Proportional to chain length









Does Teechain scale out?

Payment network deployment:

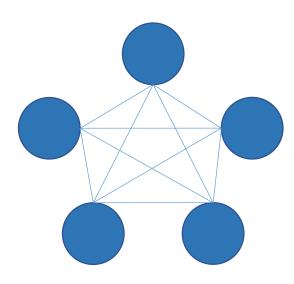
- Workload: Bitcoin transaction history across graph
- Overlay topologies: Complete vs. hub-and-spoke



Does Teechain scale out?

Payment network deployment:

Workload: Bitcoin transaction history across graph



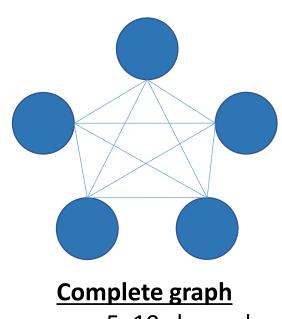
Complete graph

e.g. n=5, 10 channels (no multi-hop payments)

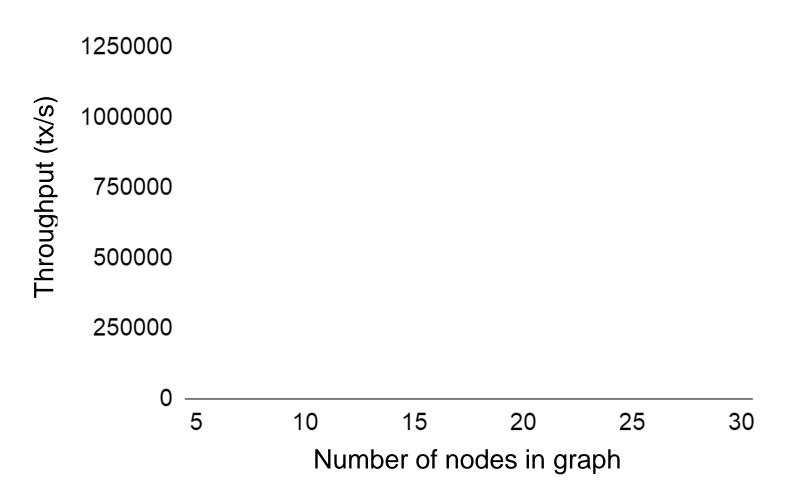


Payment network deployment:

Workload: Bitcoin transaction history across graph



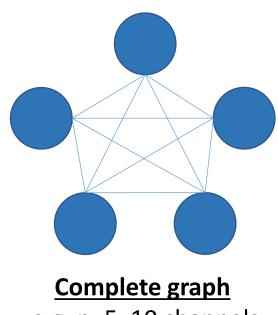
e.g. n=5, 10 channels (no multi-hop payments)

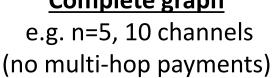


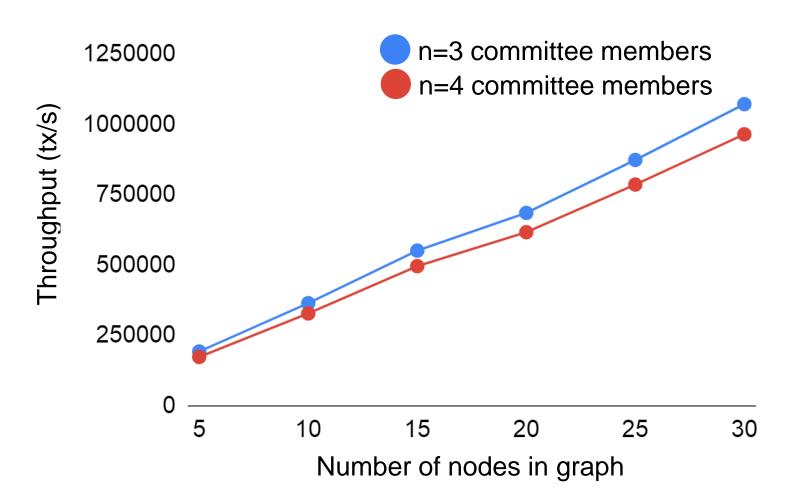


Payment network deployment:

Workload: Bitcoin transaction history across graph









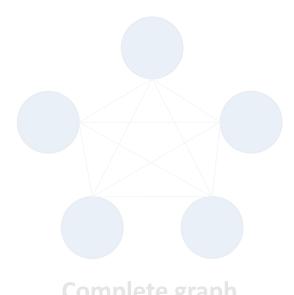
Does Teechain sc

Committee chains

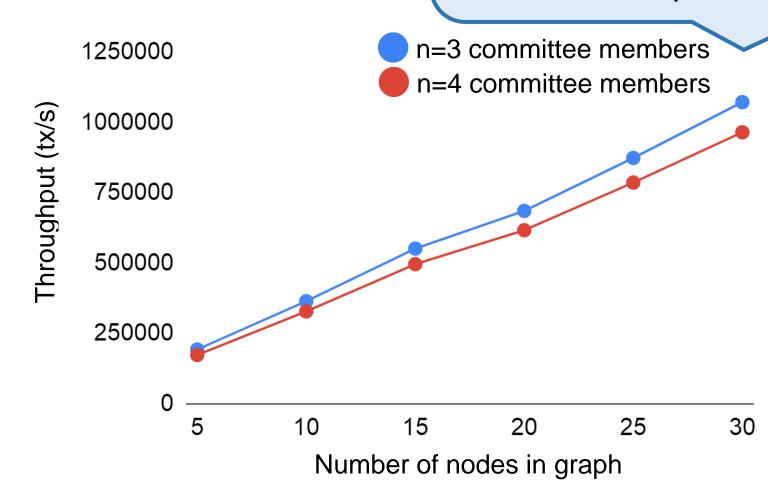
Payment network deployment

Workload: Bitcoin transaction history across graph

Throughput is limited by cost to replicate



e.g. n=5, 10 channels (no multi-hop payments)





Payment networ

- Workload: Bite

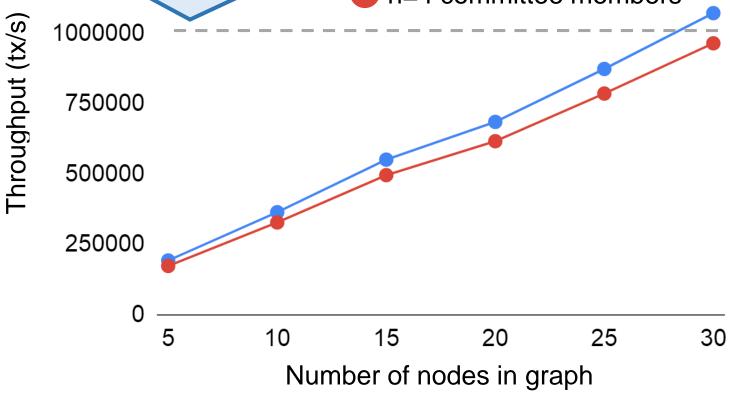
1 million tx/s

Throughput scales linearly: 30 machines

n=3 committee members n=4 committee members



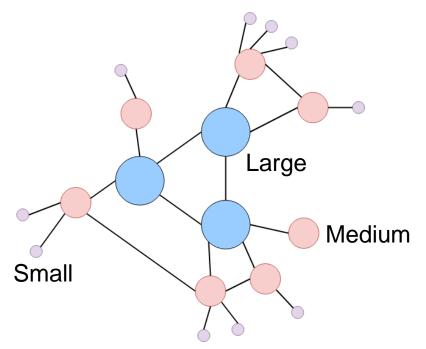
e.g. n=5, 10 channels (no multi-hop payments)





Payment network deployment:

Workload: Bitcoin transaction history across graph

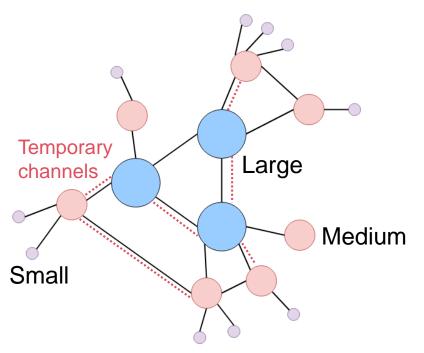


Hub-and-spoke graph



Payment network deployment:

Workload: Bitcoin transaction history across graph



Hub-and-spoke graph

Large/medium nodes use temporary channel optimization

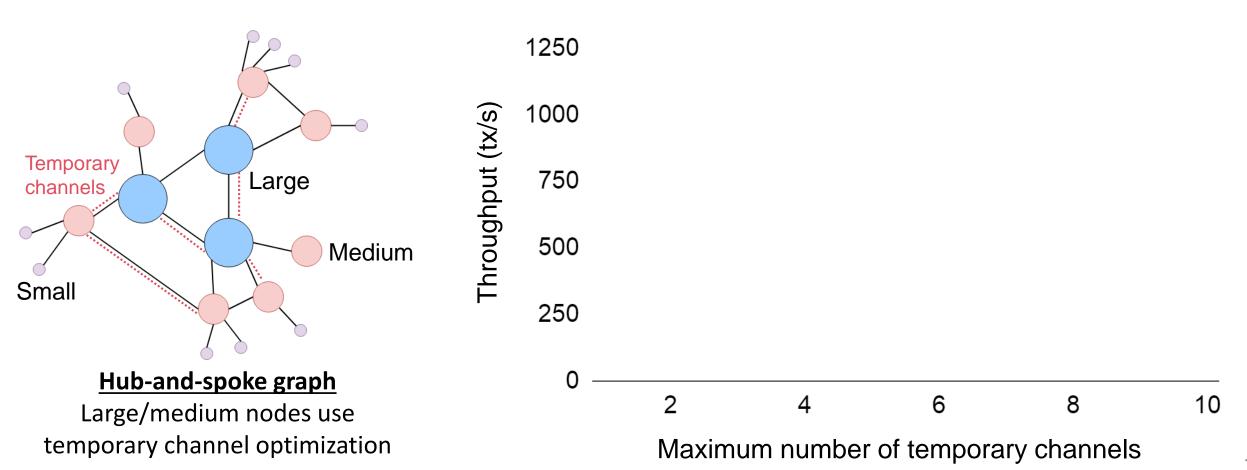
Optimization: Temporary Channels

Create temporary channels to avoid payment contention (see the paper!)



Payment network deployment:

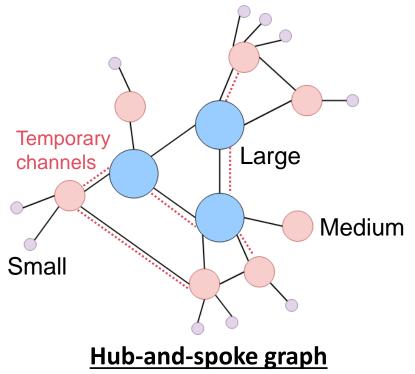
Workload: Bitcoin transaction history across graph



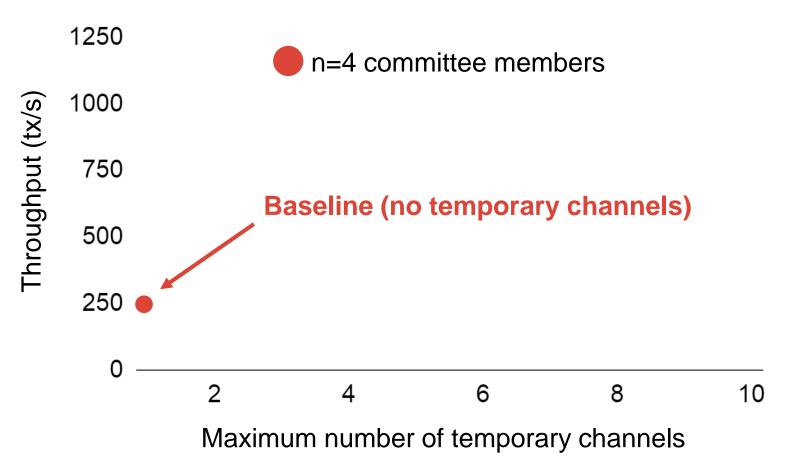


Payment network deployment:

Workload: Bitcoin transaction history across graph



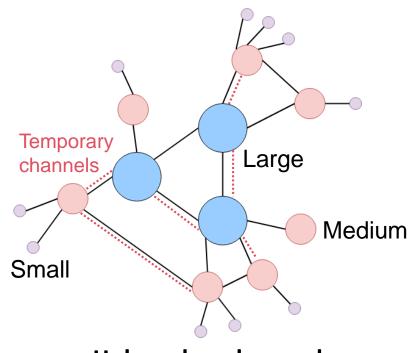
Large/medium nodes use temporary channel optimization



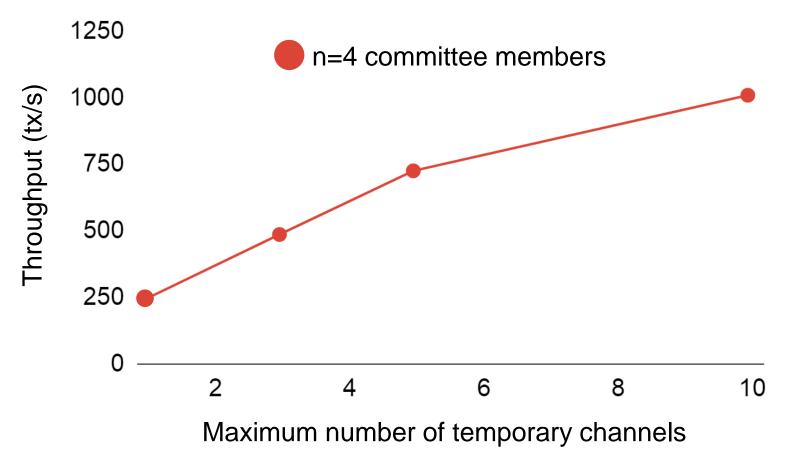


Payment network deployment:

Workload: Bitcoin transaction history across graph



Hub-and-spoke graph



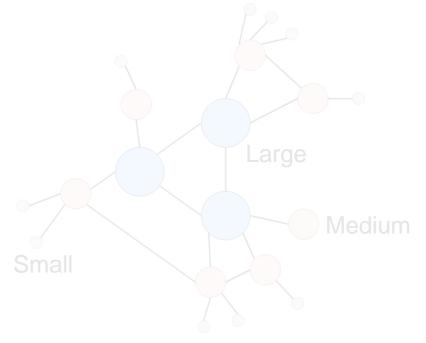


Payment network deployment:

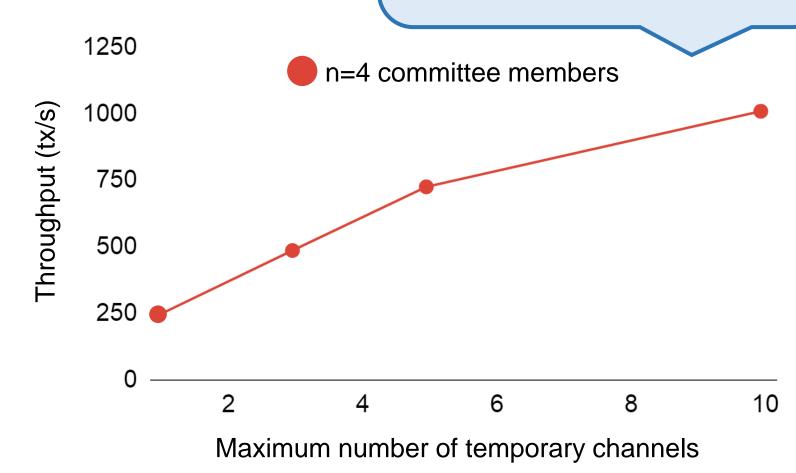
Workload: Bitcoin transaction history across graph

Optimization

Temporary channels alleviate congestion!



Hub-and-spoke graph





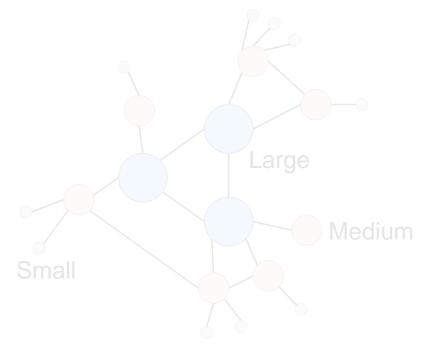
Does Teechain sca

Best performance

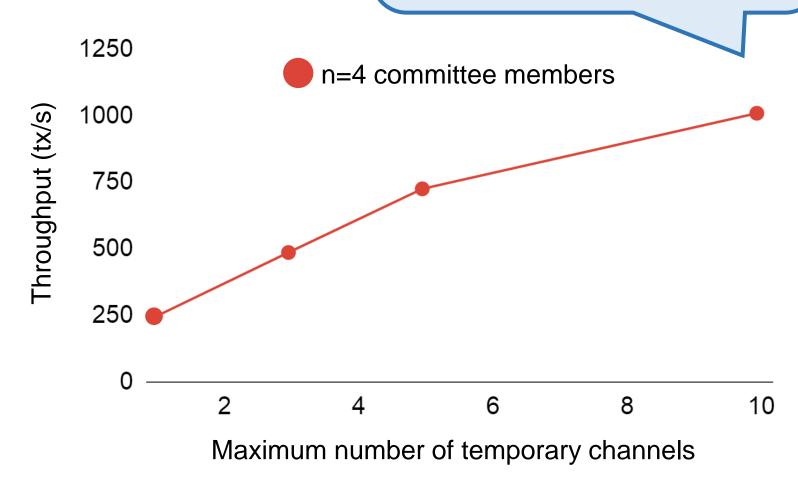
Payment network deployment:

Workload: Bitcoin transaction history across graph

Performance requires high connectivity!



Hub-and-spoke graph





Summary

Blockchains are best-effort:

- Security shouldn't rely on read/write latencies!
- Assume asynchronous blockchain access

TEEs are not silver bullets:



- Must allow for some degree of failures!
- Committees compliment TEEs

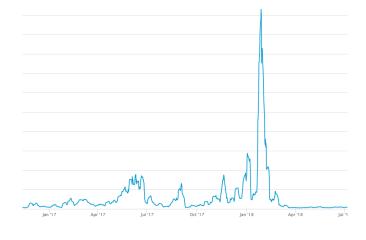
Open-source online:

- https://github.com/lsds/Teechain
- Contact us: teechain.network











Additional slides

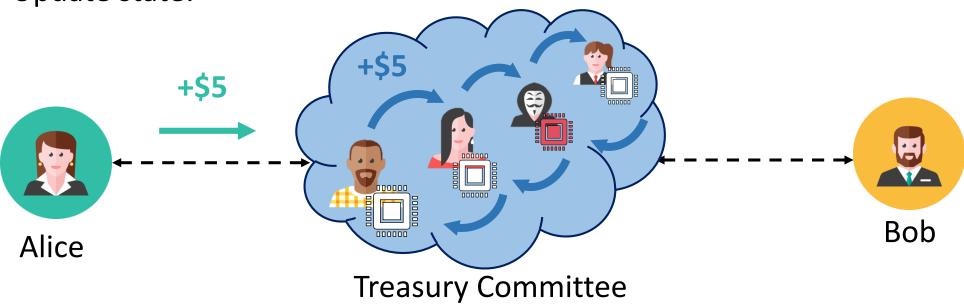


Chain replication: An overview

On each payment channel update:

- Replicate state of the head in the chain and propagate it down the chain

Update state:



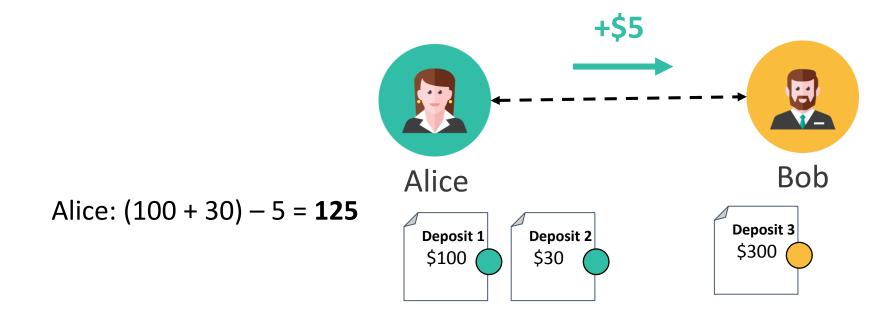


Dynamic deposits: An overview

Teechain supports dynamic deposits:

- Deposits can be added/removed from payment channels
- New deposits can be created at runtime

Collateral = Amount deposited – Amount spent





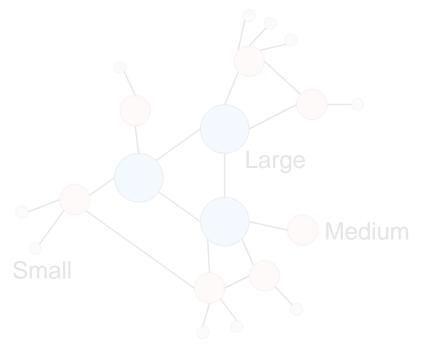
Does Teechain Sc

Throughput is Limited

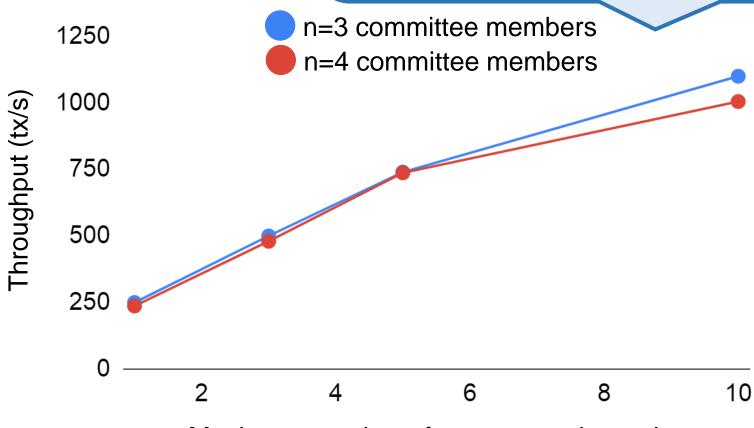
Payment Network deployment:

Workload: Bitcoin transaction history across graph

Throughput is limited by replication costs



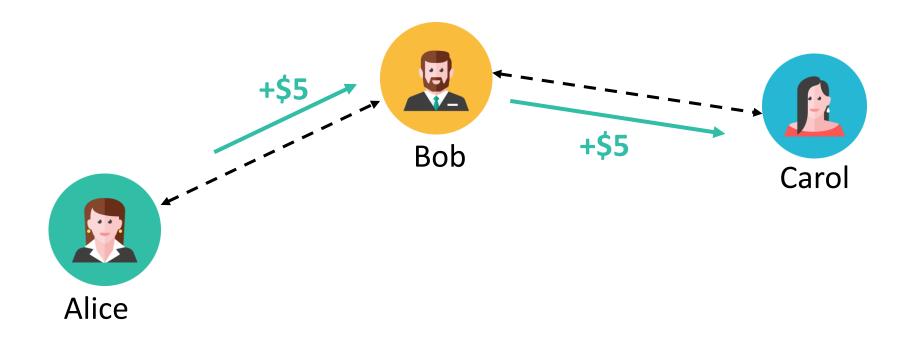
Hub-and-Spoke graph



Maximum number of temporary channels

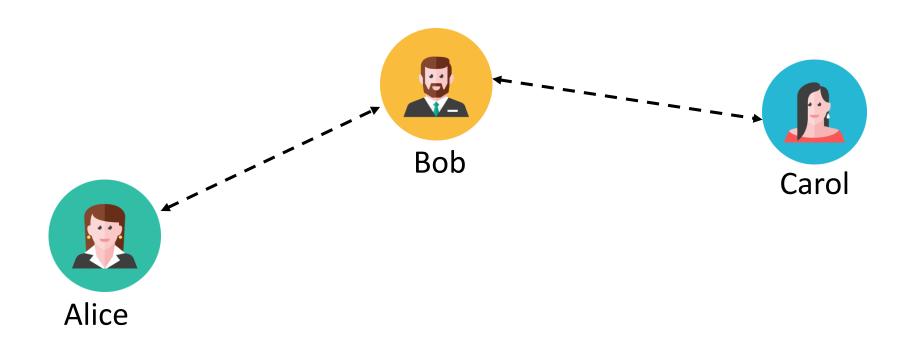


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path



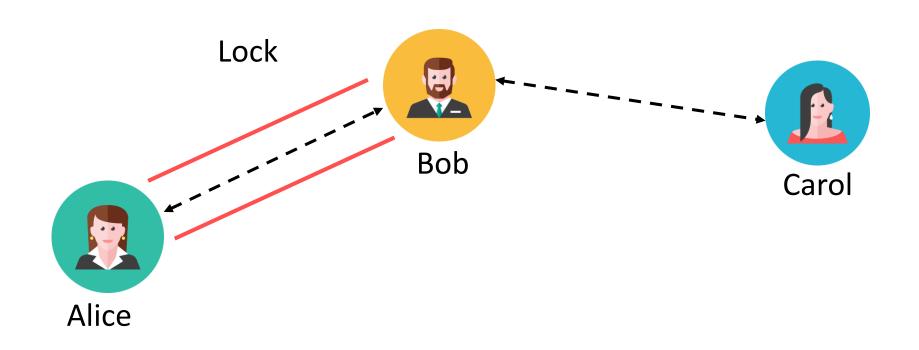


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit



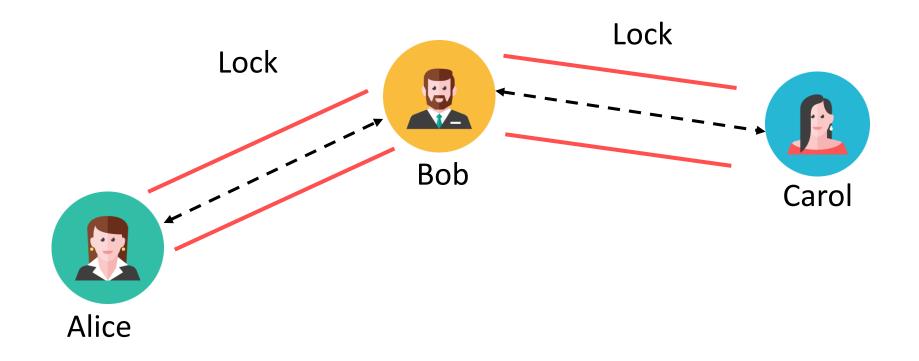


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit



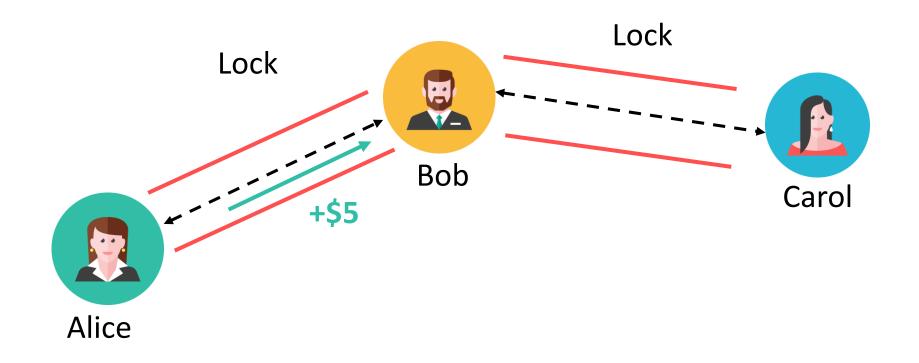


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit



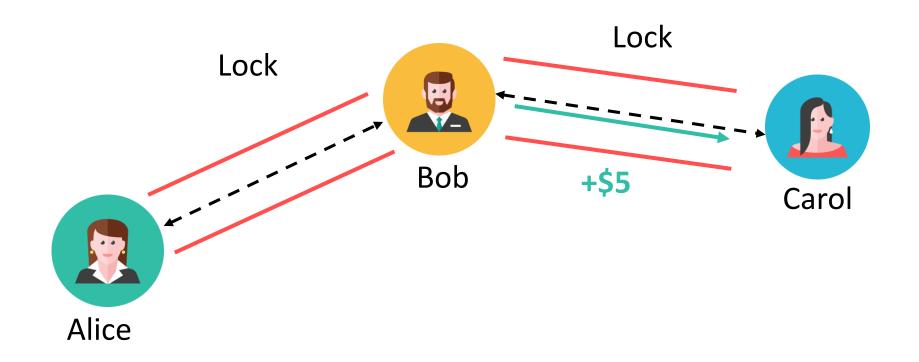


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit



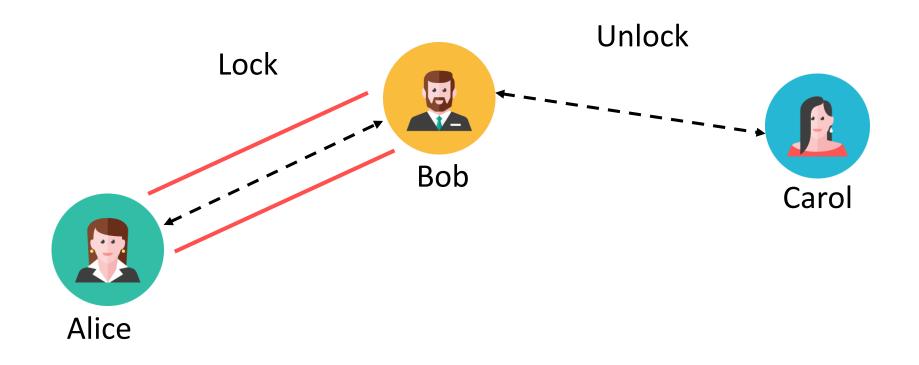


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit



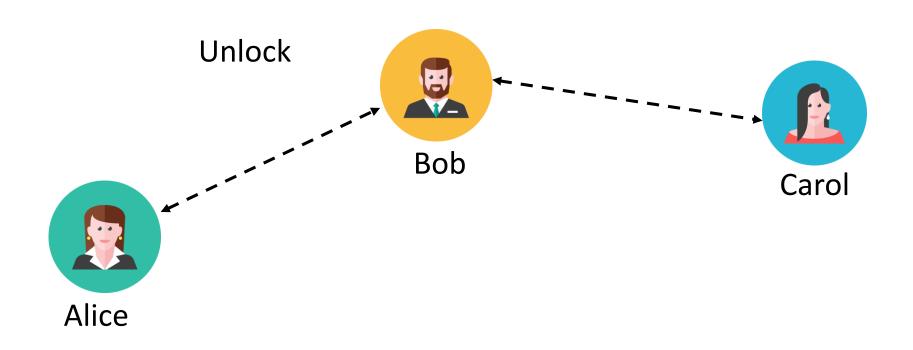


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit



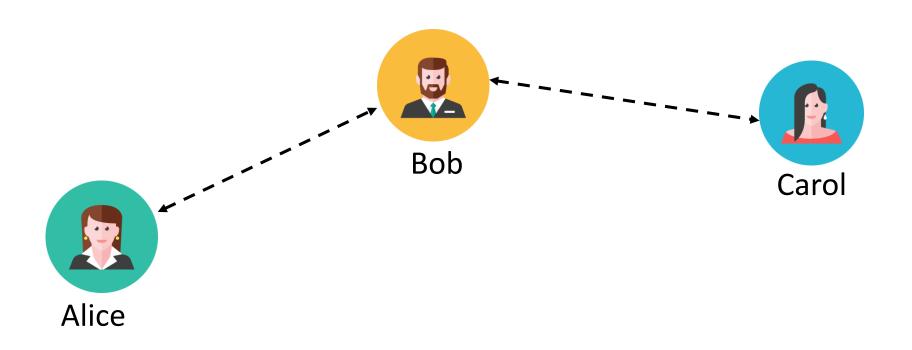


- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit





- Maintains asynchronous blockchain access
- Challenge: Ensure atomic payments across multi-hop path
- Our solution: Lock payment path and execute multi-phase commit





Μų

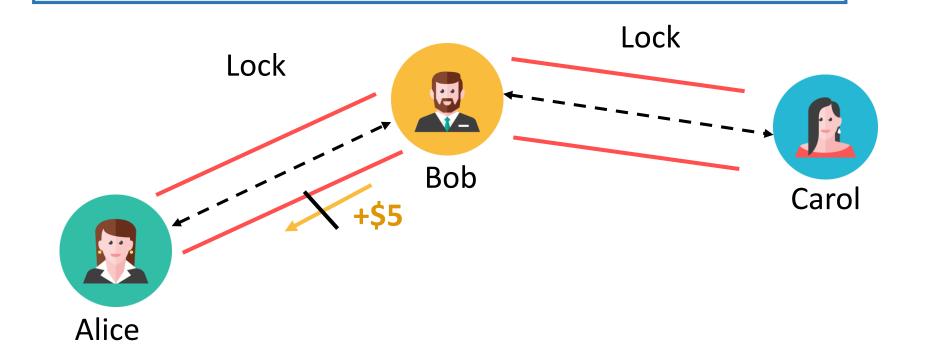
Nev

Problem: No concurrent payments!

ain access

Blocking other payments reduces throughput

ommit



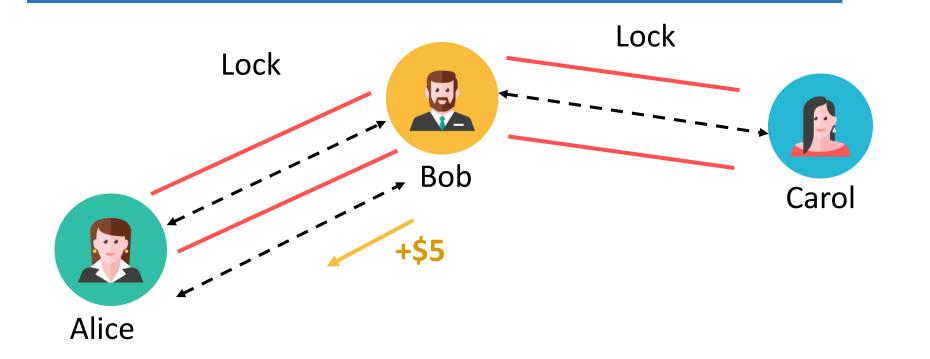


Nev

Optimization: Temporary Channels!

ain access

Dynamically create channels quickly to allow concurrent payments (see the paper!)





Nev

Optimization 2: Transaction batching!

ain access

ommit

Batch payments from the same sender to the same recipient (see the paper!)

